

# **Introduction to AI**

## **Lecture 13 Inference Rules**

**Dr. Tamal Ghosh  
Department of CSE  
Adamas University**

# Definition of Inference

In artificial intelligence, we need intelligent computers that can create new logic from old logic or by evidence, **so generating conclusions from evidence and facts is termed as Inference.**

## **Inference rules:**

Inference rules are the templates for generating valid arguments. Inference rules are applied to derive proofs in artificial intelligence, and the proof is a sequence of conclusions that leads to the desired goal.

In inference rules, the implication among all the connectives plays an important role. Following are some terminologies related to inference rules:

# Rules of Inference

- **Implication:** It is one of the logical connectives which can be represented as  $P \rightarrow Q$ . It is a Boolean expression.
- **Converse:** The converse of implication, which means the right-hand side proposition goes to the left-hand side and vice-versa. It can be written as  $Q \rightarrow P$ .
- **Contrapositive:** The negation of converse is termed as contrapositive, and it can be represented as  $\neg Q \rightarrow \neg P$ .
- **Inverse:** The negation of implication is called inverse. It can be represented as  $\neg P \rightarrow \neg Q$ .

# Truth Table from Inference Rule

- From the above term some of the compound statements are equivalent to each other, which we can prove using truth table:

<b>P</b>	<b>Q</b>	<b><math>P \rightarrow Q</math></b>	<b><math>Q \rightarrow P</math></b>	<b><math>\neg Q \rightarrow \neg P</math></b>	<b><math>\neg P \rightarrow \neg Q</math></b>
<b>T</b>	<b>T</b>	<b>T</b>	<b>T</b>	<b>T</b>	<b>T</b>
<b>T</b>	<b>F</b>	<b>F</b>	<b>T</b>	<b>F</b>	<b>T</b>
<b>F</b>	<b>T</b>	<b>T</b>	<b>F</b>	<b>T</b>	<b>F</b>
<b>F</b>	<b>F</b>	<b>T</b>	<b>T</b>	<b>T</b>	<b>T</b>

- Hence from the above truth table, we can prove that  $P \rightarrow Q$  is equivalent to  $\neg Q \rightarrow \neg P$ , and  $Q \rightarrow P$  is equivalent to  $\neg P \rightarrow \neg Q$ .

# Types of Inference rules

## 1. Modus Ponens:

The Modus Ponens rule is one of the most important rules of inference, and it states that if  $P$  and  $P \rightarrow Q$  are true, then we can infer that  $Q$  will be true. It can be represented as:

$$\text{Notation for Modus ponens: } \frac{P \rightarrow Q, P}{\therefore Q}$$

### Example:


Statement-1: "If I am sleepy then I go to bed"  $\implies P \rightarrow Q$

Statement-2: "I am sleepy"  $\implies P$

Conclusion: "I go to bed."  $\implies Q$ .

Hence, we can say that, if  $P \rightarrow Q$  is true and  $P$  is true then  $Q$  will be true.

P	Q	$P \rightarrow Q$
0	0	0
0	1	1
1	0	0
1	1	1



# Types of Inference rules

## 2. Modus Tollens:

The Modus Tollens rule states that if  $P \rightarrow Q$  is true and  $\neg Q$  is true, then  $\neg P$  will also be true. It can be represented as:

$$\text{Notation for Modus Tollens: } \frac{P \rightarrow Q, \sim Q}{\sim P}$$

**Statement-1:** "If I am sleepy then I go to bed"  $\implies P \rightarrow Q$

**Statement-2:** "I do not go to bed."  $\implies \sim Q$

**Statement-3:** Which infers that "I am not sleepy"  $\implies \sim P$

**Proof by Truth table:**

P	Q	$\sim P$	$\sim Q$	$P \rightarrow Q$
0	0	1	1	1
0	1	1	0	1
1	0	0	1	0
1	1	0	0	1

# Types of Inference rules

## 3. Hypothetical Syllogism:

The Hypothetical Syllogism rule state that if  $P \rightarrow R$  is true whenever  $P \rightarrow Q$  is true, and  $Q \rightarrow R$  is true. It can be represented as the following notation:

**Example:**

**Statement-1:** If you have my home key then you can unlock my home.  $P \rightarrow Q$

**Statement-2:** If you can unlock my home then you can take my money.  $Q \rightarrow R$

**Conclusion:** If you have my home key then you can take my money.  $P \rightarrow R$

P	Q	R	$P \rightarrow Q$	$Q \rightarrow R$	$P \rightarrow R$
0	0	0	1	1	1
0	0	1	1	1	1
0	1	0	1	0	1
0	1	1	1	1	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	0	0
1	1	1	1	1	1

# Types of Inference rules

## 4. Disjunctive Syllogism:

The Disjunctive syllogism rule state that if  $P \vee Q$  is true, and  $\neg P$  is true, then  $Q$  will be true. It can be represented as:

$$\text{Notation of Disjunctive syllogism: } \frac{P \vee Q, \neg P}{Q}$$

**Example:**

**Statement-1:** Today is Sunday or Monday.  $\implies P \vee Q$

**Statement-2:** Today is not Sunday.  $\implies \neg P$

**Conclusion:** Today is Monday.  $\implies Q$

P	Q	$\neg P$	$P \vee Q$
0	0	1	0
0	1	1	1
1	0	0	1
1	1	0	1

# Types of Inference rules

## 5. Addition:

The Addition rule is one of the common inference rule, and it states that If P is true, then PVQ will be true.

$$\text{Notation of Addition: } \frac{P}{P \vee Q}$$

### Example:

**Statement:** I have a vanilla ice-cream.  $\implies P$

**Statement-2:** I have Chocolate ice-cream.

**Conclusion:** I have vanilla or chocolate ice-cream.  $\implies (P \vee Q)$

P	Q	$P \vee Q$
0	0	0
1	0	1
0	1	1
1	1	1


# Types of Inference rules

## 6. Simplification:

The simplification rule state that if  $P \wedge Q$  is true, then  $Q$  or  $P$  will also be true. It can be represented as:

$$\text{Notation of Simplification rule: } \frac{P \wedge Q}{Q} \text{ Or } \frac{P \wedge Q}{P}$$

P	Q	$P \wedge Q$
0	0	0
1	0	0
0	1	0
1	1	1



# Types of Inference Rules

## 7. Resolution:

The Resolution rule states that if  $P \vee Q$  and  $\neg P \wedge R$  is true, then  $Q \vee R$  will also be true. It can be represented as

$$\text{Notation of Resolution} = \frac{A \vee B, \neg A \vee C}{B \vee C}$$

A	B	$\neg B$	C	$A \vee B$	$\neg B \vee C$	$A \vee C$
1	0	1	1	1	1	1
1	1	0	1	1	1	1
0	1	0	1	1	1	1
0	0	1	1	0	1	1
1	0	1	0	1	1	1
1	1	0	0	1	0	1
0	1	0	0	1	0	0
0	0	1	0	0	1	0